Research Project Proposal: Multi-Agent Path Finding with Removable Obstacles

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Introduction

- Multi-Agent Path Finding (MAPF) is a challenging Artificial Intelligence topic that has many applications
- The task in MAPF problems is to find non-interfering paths for multiple agents, each one with a unique start and goal position in a known environment





Applications



Applications





Applications



Preliminaries







- Simplifying assumptions
 - Point robots
 - No kinematic constraints
 - Discretized environment





Preliminaries

- In the MAPF problem we are given:
 - A graph G = (V, E)
 - A set of k labeled agents $A = \{a_1, \dots, a_k\}$
 - A set of k start positions $s = \{s_1, \dots, s_k\}, s_i \in V$
 - A set of k goal positions $g = \{g_1, \dots, g_k\}, g_i \in V$
- Goal: Find a set of k non-conflicting paths, one for each agent, that minimize some objective function
 - Makespan: latest arrival time of an agent at its goal location
 - Flowtime: sum of the arrival times of all the agents at their goal locations

Preliminaries

- A vertex conflict is a tuple (a_1, a_2, v, t) meaning that agents a_1 and a_2 are occupying the same vertex v at timestamp t
- An edge conflict is a tuple (a_1, a_2, v, v', t) meaning that, from timestamp t to t + 1, agent a_1 is traveling from v to v' while a_2 is traveling from v' to v
- which brings a_i from $v_0 = s_i$ to $v_m = g_i$
- A valid solution is defined as a set of k conflict-free paths (no collisions)

• Time is discretized into timesteps and, at each timestep, every agent can either change location moving to an adjacent vertex or wait at its current position

 \mathcal{V}







• A path π_i for an agent a_i can be modeled as a sequence of vertices (v_0, \dots, v_m)

NP-hardness

- for both makespan and flowtime minimization



• Theorem [Yu and LaValle, 2013]: MAPF problems are NP-hard to solve optimally

NP-hardness proof is based on a direct reduction from the 3-satisfiability problem

Optimal MAPF Solvers

Optimal MAPF solvers

Search based

A* [Hart et al., 1968]
M* [Wagner and Choset, 2011]
ICTS [Sharon et al., 2011]
CBS [Sharon et al., 2012]

Reduction based

SAT [Surynek, 2012] **ASP** [Erdem et al., 2013] **ILP** [Yu and LaValle, 2013]

Conflict-Based Search

- Two-level structure: high and low level
- account constraints imposed by the high-level
- and imposes new constraints to the agents in order to avoid the conflict
 - The idea is to explore all possible ways to solve the conflict

• The low-level searches for an optimal path for each agent individually taking into

• When a conflict between individual paths is found, the high-level expands a constraint tree via a split action, keeping all the information about the conflict,

• A constraint prohibits an agent from being in a certain position at a given time

Conflict-Based Search

The red and blue agents collide in the yellow cell (x=3,y=2) at time 1

Add constraint: the red agent is not allowed to be in cell (3,2) at time 1





Add constraint: the blue agent is not allowed to be in cell (3,2) at time 1



Conflict-Based Search

- We have just seen the general idea behind CBS-based solvers
- An improved version of CBS is available with the name of Improved Conflict-Based Search (ICBS) [Boyarski et al., 2015]
 - ICBS-h (or CBSH) [Felner et al., 2018] is the state-of-the-art CBS-based solver

Non-static environments

- In classical MAPF problems, we make strong assumptions concerning the environment
- Real-world environments may have different features and properties
 - Dynamic obstacles
 - Destructible obstacles
 - Shrinking maps
 - Toll booths



Reconfigurable environments

- We want to introduce the concept of *environment reconfiguration*
 - The idea is that something in the environment can be changed, at will, under certain conditions
 - In general, an environment that is mutable in this sense is called a reconfigurable environment
- Our first goal is to introduce and clearly define the concept of environment reconfiguration in the MAPF context

Environments with removable obstacles







MAPF with removable obstacles

- assume that some obstacles can just be removed from the environment
- problem could be extended to arbitrary graphs
- A removable obstacle is defined as a set of (blocked) cells and each removable obstacle has an associated cost called removal cost
 - We assume that there are no cells shared by multiple obstacles
- We are given a budget that can be spent on removing obstacles

• We focus on reconfigurable environments with destructible obstacles, so we

• For simplicity we focus on grids, composed by free and blocked cells, but the



MAPF with removable obstacles

- paths without any collision
- A unitary budget and two removable obstacles:
 - λ_1 composed by two cells and a unitary cost
 - λ_2 composed by one cell and a unitary cost



• For optimally solving a MAPF with removable obstacles (RO-MAPF) problem we mean returning the solution with the minimum cost and a set of removable obstacles to remove (within the budget) in order to allow agents to follow their

(b)

MAPF with removable obstacles

- MAPF problem is NP-hard to solve optimally, even on grid-based environments
- RO-MAPF problem is NP-hard to solve optimally too since it is a generalization of the MAPF problem
- Despite this intractability issue, our goal is to develop an effective CBS-based algorithm able to solve RO-MAPF problems optimally
 - For now, we have designed a draft version of our final solver

Draft version of our RO-MAPF solver



Research timeline



- We are keeping in mind some conference deadlines
 - AAAI 2020: deadline towards the beginning of September
 - AAMAS 2020 and ICAPS 2020: deadline in mid-November



For more information about MAPF see the *brand new* website <u>http://mapf.info</u>